



LEGALM: Efficient Legalization for Mixed-Cell-Height Circuits with Linearized Augmented Lagrangian Method

Jing Mai¹, Chuanyuan Zhao¹, Zuodong Zhang¹, Zhixiong Di², Yibo Lin¹, Runsheng Wang¹, Ru Huang¹

¹Peking University

²Southeast Jiaotong University

Mar 17, 2025



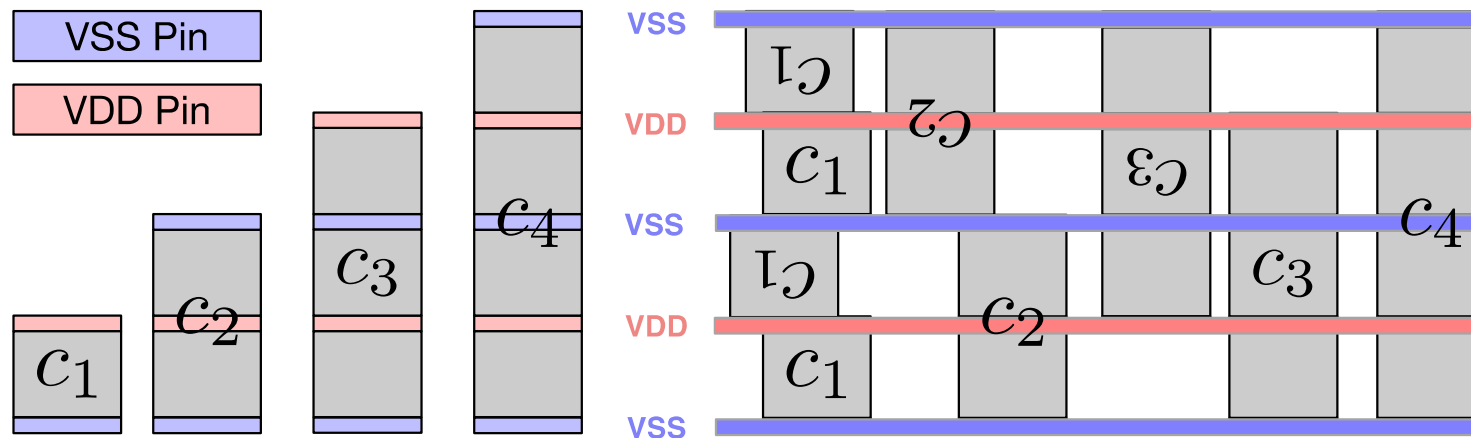
1. Introduction
2. The LEGALM Algorithm
3. Experimental Results
4. Conclusion

Introduction

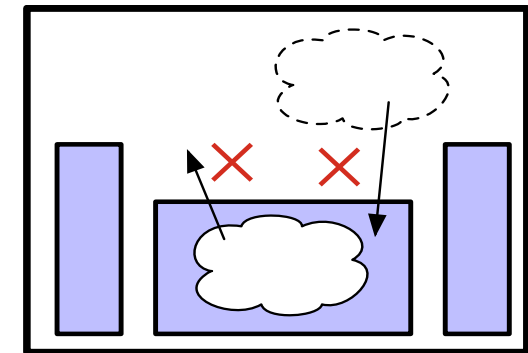
Mixed-Cell-Height Circuits & Fence Regions



- Recent mixed-cell-height designs combine higher and smaller cells to optimize PPA in modern ASICs.
 - Higher cells enhance performance and routability for critical paths.
 - Smaller cells improve area efficiency and reduce power for non-critical logic.
- Modern ASIC CAD tools provide the fence region as an important feature.



Fence Region

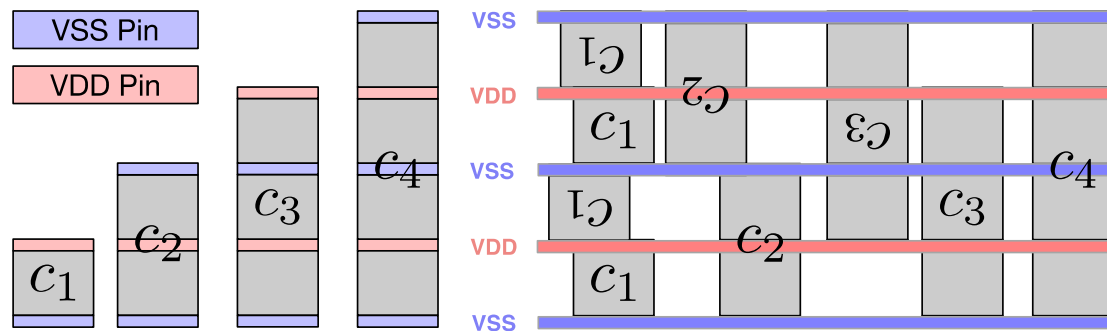


Challenges & Motivation



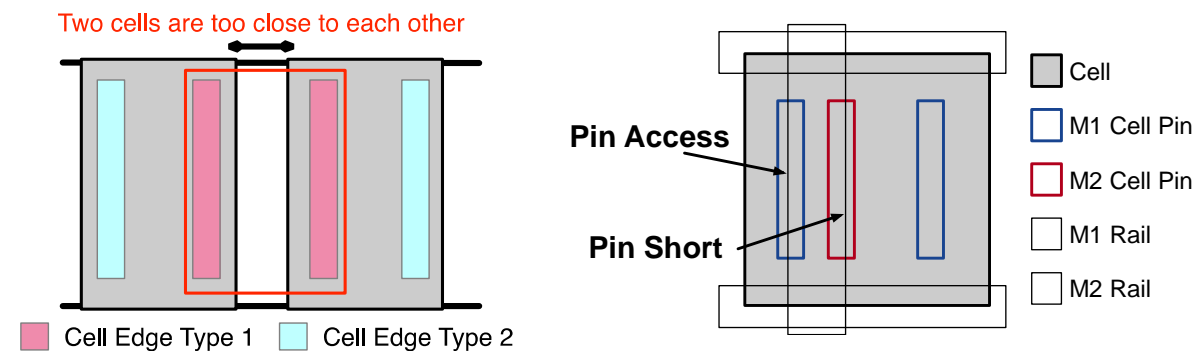
□ Why Legalization matters:

- ▶ Eliminates design rule violations post-global placement.
- ▶ Impacts downstream routing and performance.



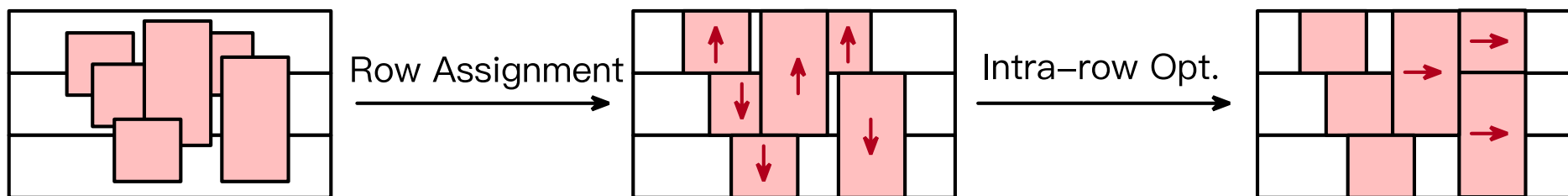
□ Challenges in Mixed-Cell-Height:

- ▶ Site alignment, overlap-free.
- ▶ Cross-row shapes, P/G alignment, fence regions, edge spacing, pin access.



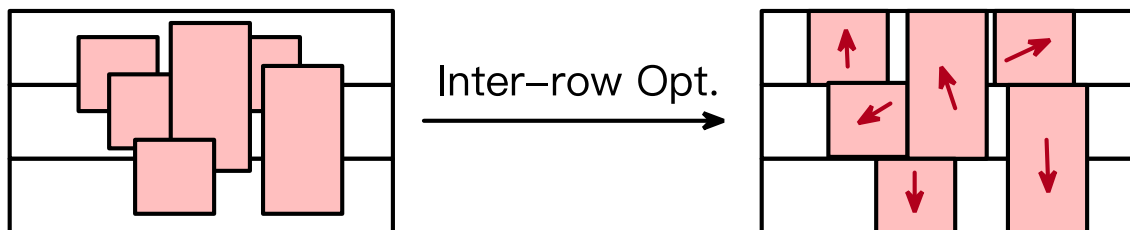
□ Intra-row Methods:

- ▶ Row assignment + row-bases optimization algorithm (e.g., ILP [Li+,DAC'18], LCP [Chen+, DAC'17]).
- ▶ Limited vertical movement space during optimization.



□ Inter-row Methods:

- ▶ Network flow / ILP-based (e.g., [Darav+, ISPD'17])
- ▶ High flexibility but computationally expensive.

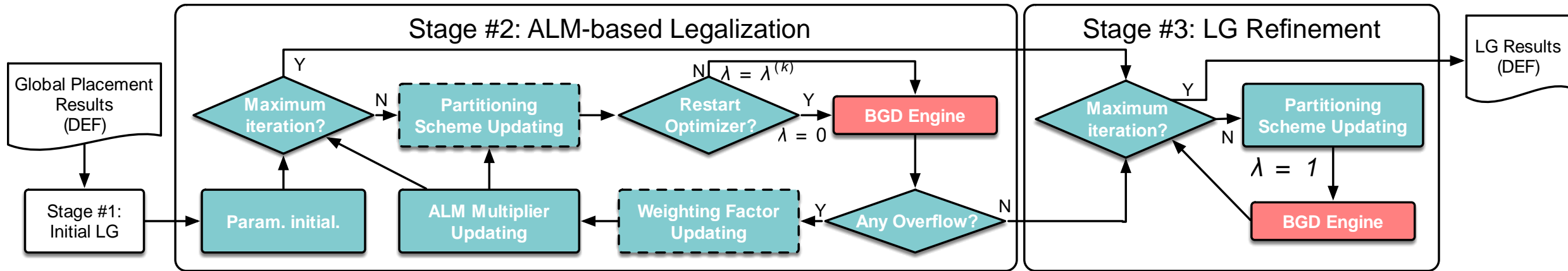


LEGALM, an efficient inter-row legalization method for mixed-cell-height circuits with fence region constraints using the linearized augmented Lagrangian method.

1. **Augmented Lagrangian Method (ALM)** for efficient vertical and horizontal cell movements.
2. **Block Gradient Descent (BGD)** for parallel cell updates.
3. **Triple-fold partitioning** for GPU efficiency.

Results: 6-36% better quality, $2.25\text{-}5.99\times$ speedup on million-cell designs.

The LEGALM Algorithm



❑ 3-Stage Workflow:

1. **Initial Legalization:** Minimal displacement, ignore overlaps.
2. **ALM-based Legalization:** Optimize displacement + eliminate overlaps.
3. **Refinement:** Strict no-overlap optimization.

❑ Key Components:

- Linearized ALM formulation for constraint relaxation.
- BGD for parallel cell updates.
- GPU-friendly partitioning.

Legalization Formulation



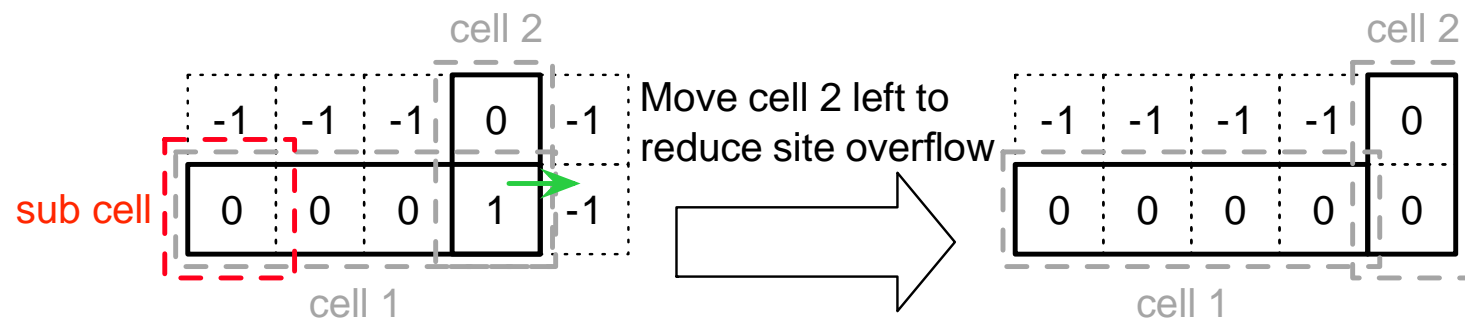
$$\min \sum_{j \in S} \sum_{i \in N} \sum_{t \in \mathcal{T}_i} w_{i,t,j} x_{i,t,j},$$

$$s.t. x_{i,t,j} \in \{0, 1\},$$

$$g_j(\mathbf{x}) = \sum_{i \in N} \sum_{t \in \mathcal{T}_i} x_{i,t,j} - 1 \leq 0, \quad \forall j \in S,$$

connected sub-cell constraints,

fence region constraints,



- \mathcal{T}_i Partition cell i into sub-cells of width 1 and height H (row height).
- $x_{i,t,j}$ Binary indicator if sub-cell t ($t \in \mathcal{T}_i$) of cell i is placed at site j ($j \in S$).
- $w_{i,t,j}$ Displacement cost for placing sub-cell t ($t \in \mathcal{T}_i$) at site j ($j \in S$).
- $g_j(\mathbf{x})$ Overflow at site j ($j \in S$).
- **Objective:** Minimize displacement while satisfying constraints.

Augmented Lagrangian Formulation



- **Lagrangian relaxation** to handle overlap-free constraints.
- **Slack variables** r_j for site overflow.

$$\begin{aligned}\mathcal{L}(\mathbf{x}, \mathbf{r}, \boldsymbol{\lambda}) = & \sum_{j \in S} \sum_{i \in N} \sum_{t \in \mathcal{T}_i} w_{i,t,j} x_{i,t,j} \\ & + \sum_{j \in S} \lambda_j \left[(g_j(\mathbf{x}) + r_j) + \frac{\sigma}{2} (g_j(\mathbf{x}) + r_j)^2 \right] \\ & + I_{\mathcal{X}}(\mathbf{x}),\end{aligned}$$

- Solve for optimal r_j by **elimination method**: $r_j = \max(0, -\frac{1}{\sigma} - g_j(\mathbf{x}))$.
- Iterative multiplier updates based on **KKT conditions**.

$$\text{Sub-problem: } \mathbf{x}^{(k+1)} = \underset{\mathbf{x}}{\operatorname{argmin}} \mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) = \underset{\mathbf{x}}{\operatorname{argmin}} \psi(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) + I_{\mathcal{X}}(\mathbf{x}).$$

$$\lambda_j^{(k+1)} = \max \left(\lambda_j^{(k)} + h_f \cdot \left[(g_j(\mathbf{x}) + r_j) + \frac{\sigma}{2} (g_j(\mathbf{x}) + r_j)^2 \right], 0 \right),$$

Augmented Lagrangian Formulation



- **Lagrangian relaxation** to handle overlap-free constraints.
- **Slack variables** r_j for site overflow.

$$\begin{aligned}\mathcal{L}(\mathbf{x}, \mathbf{r}, \boldsymbol{\lambda}) = & \sum_{j \in S} \sum_{i \in N} \sum_{t \in \mathcal{T}_i} w_{i,t,j} x_{i,t,j} \\ & + \sum_{j \in S} \lambda_j \left[(g_j(\mathbf{x}) + r_j) + \frac{\sigma}{2} \boxed{(g_j(\mathbf{x}) + r_j)^2} \right] \quad \text{Variable Coupling} \\ & + I_{\mathcal{X}}(\mathbf{x}),\end{aligned}$$

$(x_1 + x_2 + \dots + x_n)(x_1 + x_2 + \dots + x_n)$

- Solve for optimal r_j by **elimination method**: $r_j = \max(0, -\frac{1}{\sigma} - g_j(\mathbf{x}))$.
- Iterative multiplier updates based on **KKT conditions**.

$$\text{Sub-problem: } \mathbf{x}^{(k+1)} = \underset{\mathbf{x}}{\operatorname{argmin}} \mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) = \underset{\mathbf{x}}{\operatorname{argmin}} \psi(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) + I_{\mathcal{X}}(\mathbf{x}).$$

$$\lambda_j^{(k+1)} = \max \left(\lambda_j^{(k)} + h_f \cdot \left[(g_j(\mathbf{x}) + r_j) + \frac{\sigma}{2} (g_j(\mathbf{x}) + r_j)^2 \right], 0 \right),$$

Sub-problem: $\mathbf{x}^{(k+1)} = \underset{\mathbf{x}}{\operatorname{argmin}} \mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) = \underset{\mathbf{x}}{\operatorname{argmin}} \psi(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) + I_{\mathcal{X}}(\mathbf{x}).$

- Given fixed multiplier $\boldsymbol{\lambda}$, let $f(\mathbf{x}) = \psi(\mathbf{x}, \boldsymbol{\lambda})$, solve the following function

$$\min_{\mathbf{x}} f(\mathbf{x}) + I_{\mathcal{X}}(\mathbf{x}).$$

- Proximal Mapping Operator** (can be proved a linear operator).

$$P_{\mathcal{X}}(\mathbf{v}) = \arg \min_{\mathbf{u} \in \mathcal{X}} \|\mathbf{u} - \mathbf{v}\|_2^2,$$

- Proximal Gradient Method (diamond search)

$$\mathbf{x}^{(k+1)} \in P_{\mathcal{X}} \left(\mathbf{x}^{(k)} - \tau \nabla f(\mathbf{x}^{(k)}) \right),$$

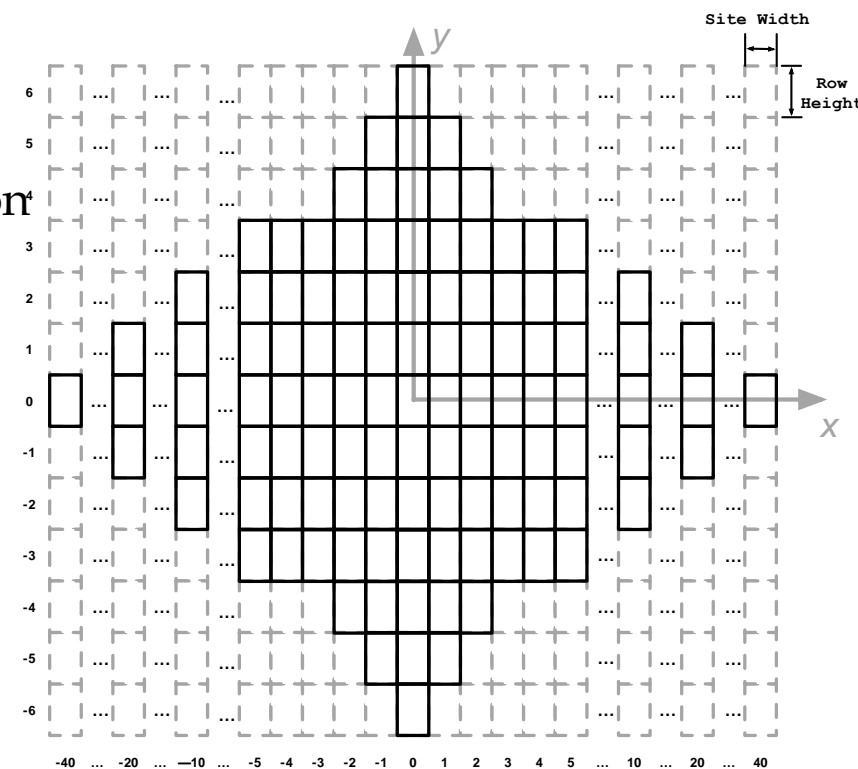


Figure: Diamond search range for the cell located at the red site.

□ Final Cost Function:

$$cost_{i,t,j} = w_{i,t,j} + \lambda_j \max\{1 + \sigma g_j(\mathbf{x}), 0\} + p \cdot H \cdot R_{i,t,j}$$

Displacement Cost

Overflow Cost

Routability Penalty

□ Steps:

- Enumerate candidate positions (diamond search).
- Compute costs for displacement, overflow and technology.
- Select minimum-cost positions.

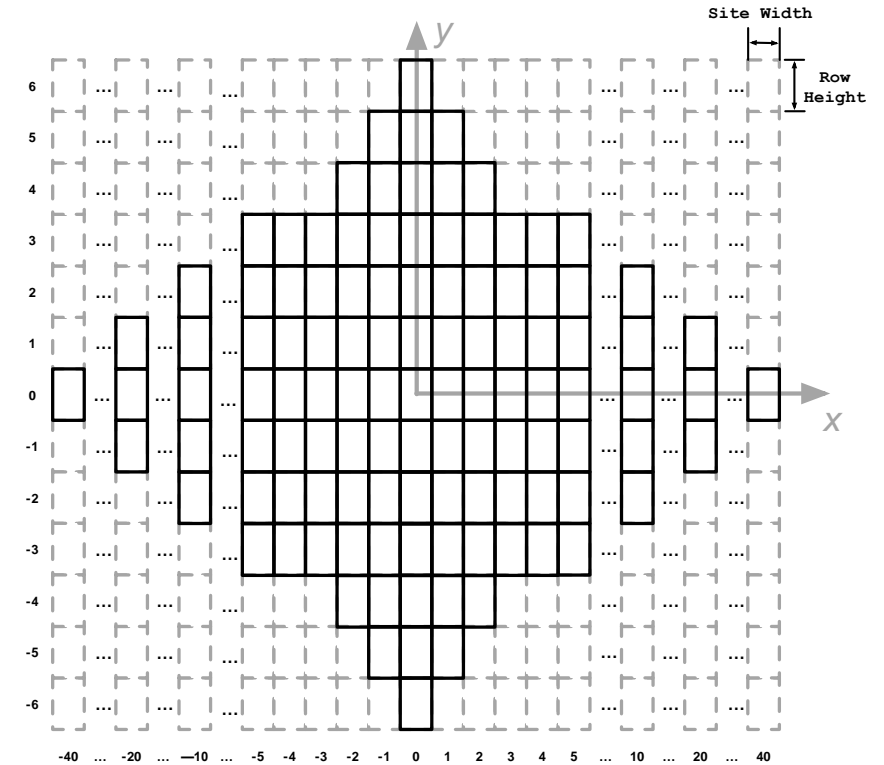


Figure: Diamond search range for the cell located at the red site.

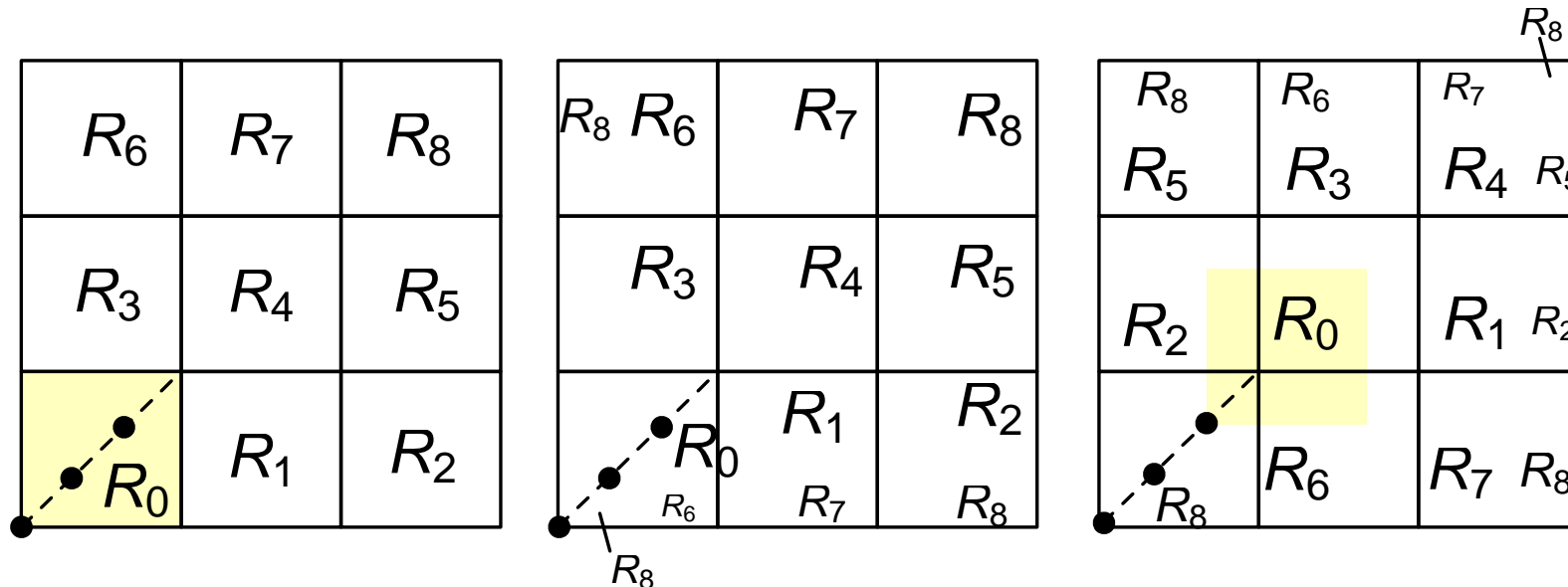
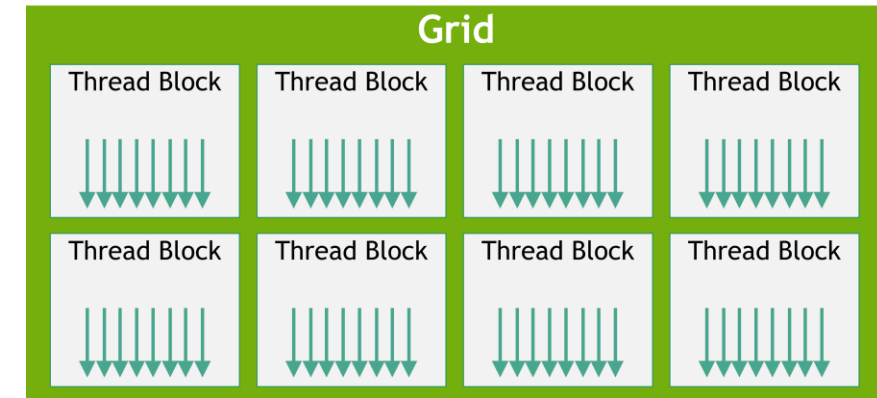
Update single cell \rightarrow Update distant cells in parallel

□ Triple-fold Partitioning:

- Divide layout into non-overlapping sub-regions (totally 3 partition schemes).
- Rotate partitions schemes to avoid update conflicts.

□ GPU Acceleration:

- Kernel design: 1 grid per sub region, threads for candidate selections.
- Shared memory for cost reduction.



Experimental Results

□ Benchmarks:

- ICCAD-2017 (routability / fence regions). [Darav+, ICCAD'17]
- Modified ISPD-2015 (million-cell designs). [Chow+, DAC'16]

□ Metrics:

- Quality score \mathcal{S} : Combine HPWL variation \mathcal{S}_{hpwl} , weighted averaged displacement \mathcal{S}_{am} , maximum displacement \mathcal{M}_{max} , #pin short/access DRVs N_p , and #edge spacing DRVs N_e .
- Runtime and speedup.

□ Platform:

- One NVIDIA A800 GPU.
- Two Intel Xeon Platinum 8358 CPUs (2.60GHz, 32 cores) with 1024GB RAM.

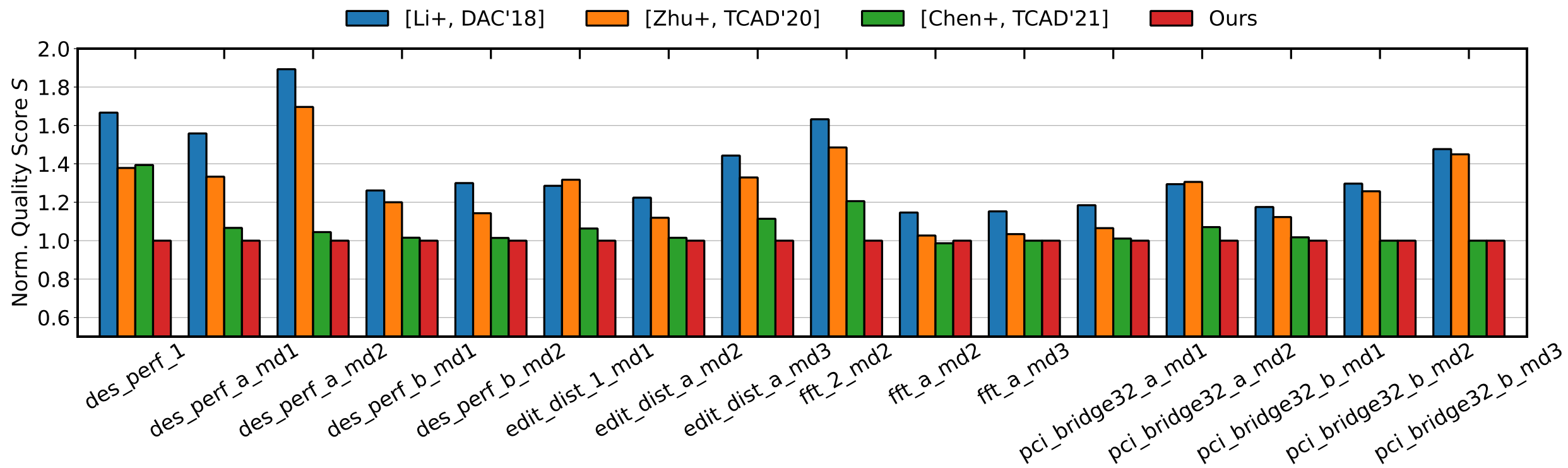
Table: Statistics of ICCAD-2017 Benchmarks.

Case	#Cells of Different Heights (H)				Den. (%)	#Regions
	1	2	3	4		
des_perf_1	112644	0	0	0	90.6	0
des_perf_a_md1	103589	4699	0	0	55.1	4
des_perf_a_md2	105030	1086	1086	1086	55.9	4
des_perf_b_md1	106782	5862	0	0	55.0	12
des_perf_b_md2	101908	6781	2260	1695	64.7	12
edit_dist_1_md1	118005	7994	2664	1998	67.4	0
edit_dist_a_md2	115066	7799	2599	1949	59.4	1
edit_dist_a_md3	119616	2599	2599	2599	57.2	1
fft_2_md2	28930	2117	705	529	82.7	0
fft_a_md2	27431	2018	672	504	32.3	0
fft_a_md3	28609	672	672	672	31.2	0
pci_bridge32_a_md1	26680	1792	597	448	49.5	4
pci_bridge32_a_md2	25239	2090	1194	994	57.7	4
pci_bridge32_b_md1	26134	1756	585	439	26.6	3
pci_bridge32_b_md2	28038	292	292	292	18.3	3
pci_bridge32_b_md3	27452	292	585	585	22.2	3

Table: Statistics of ISPD-2015 Benchmarks.

Case	#Cells of Different Heights (H)		Den. (%)
	1	2	
mgc_superblue11_a	861314	64302	43
mgc_superblue12	1172586	114362	45
mgc_superblue14	564769	47474	56
mgc_superblue16_a	625419	47474	48
mgc_superblue19	478109	27988	52

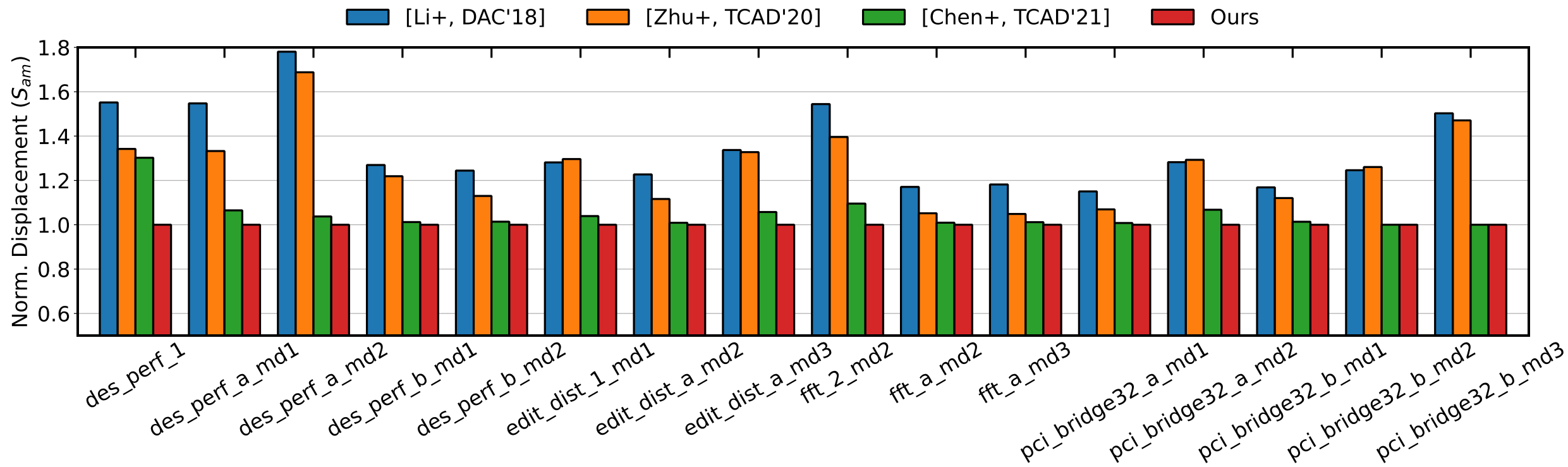
- **Quality:** LEGALM improve overall score S by 6-36% vs. SOTA with $1.03\text{-}3.83\times$ speedup.



Result on ICCAD-2017 (cont'd)



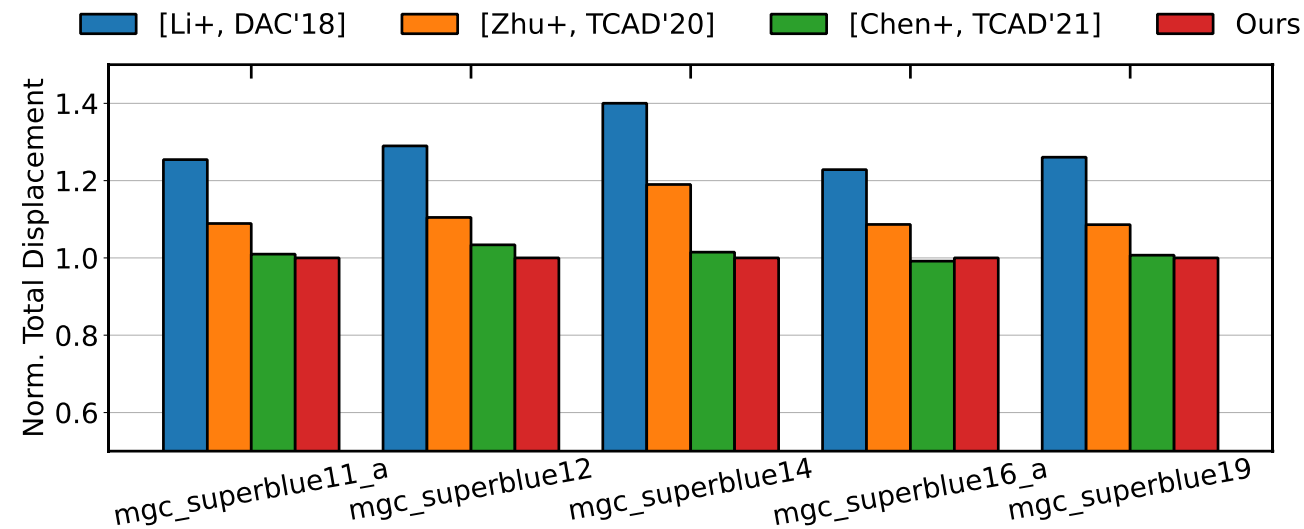
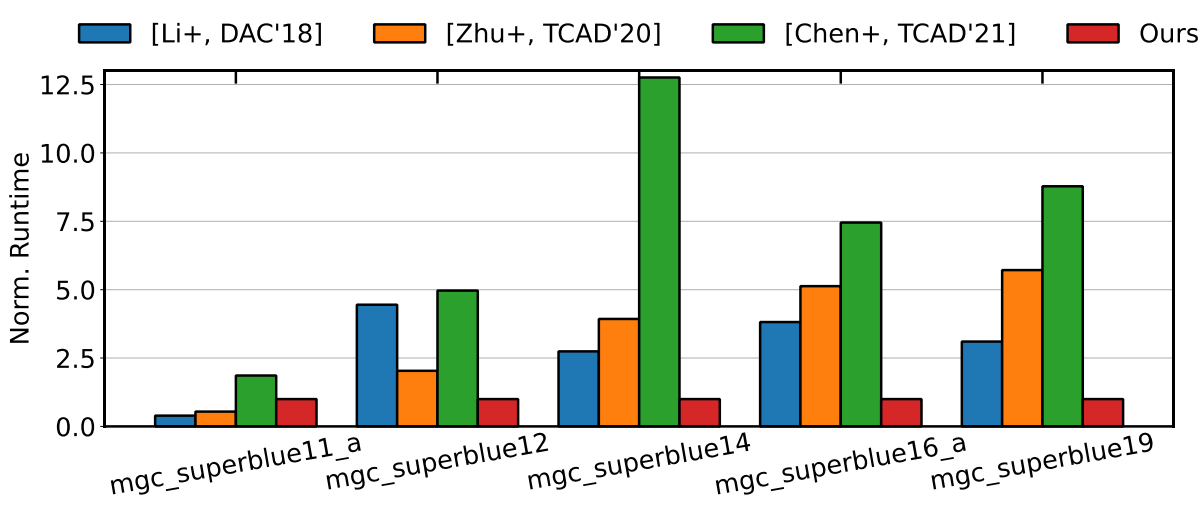
- **Quality:** LEGALM improve overall score S by 6-36% vs. SOTA with $1.03\text{-}3.83\times$ speedup.
- **Breakdown:**
 - 13-34% better HPWL variation S_{hpwl} .
 - 4.4-33.2% better weighted average displacement S_{am} .
 - 10% better maximum displacement M_{max} .



Scalability on Large Designs



- **Cases:** mgc_superblue (1M+ cells)
- **Results:**
 - 2.25–5.99× faster than SOTA.
 - 1.1–28.5% better displacement than SOTA.
 - Legalization in < 10 seconds for 3/5 cases.
- **Why?:** GPU parallelism + efficient partitioning.



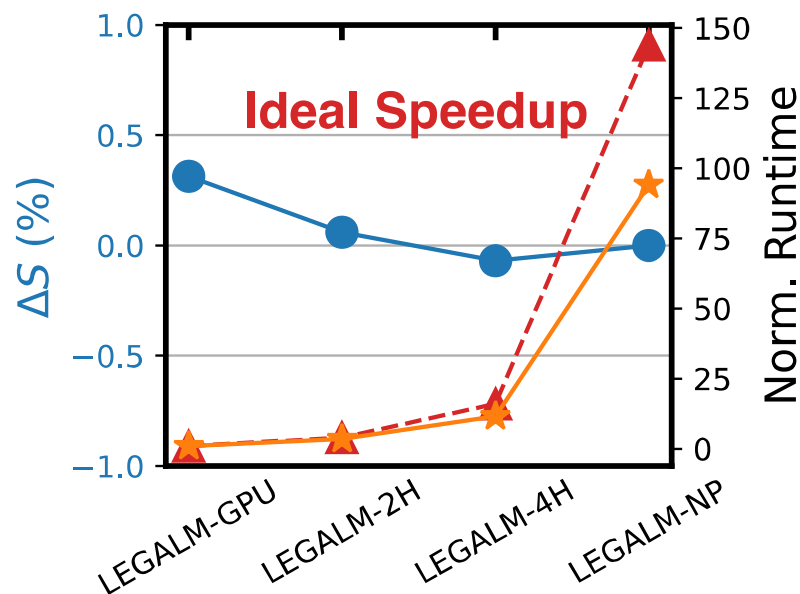
Ablation Study & Runtime Breakdown



□ Triple-fold Partitioning (TP) Impact:

- 94.2 \times speedup vs. no partitioning (LEGALM-NP).
- <0.5% quality loss.

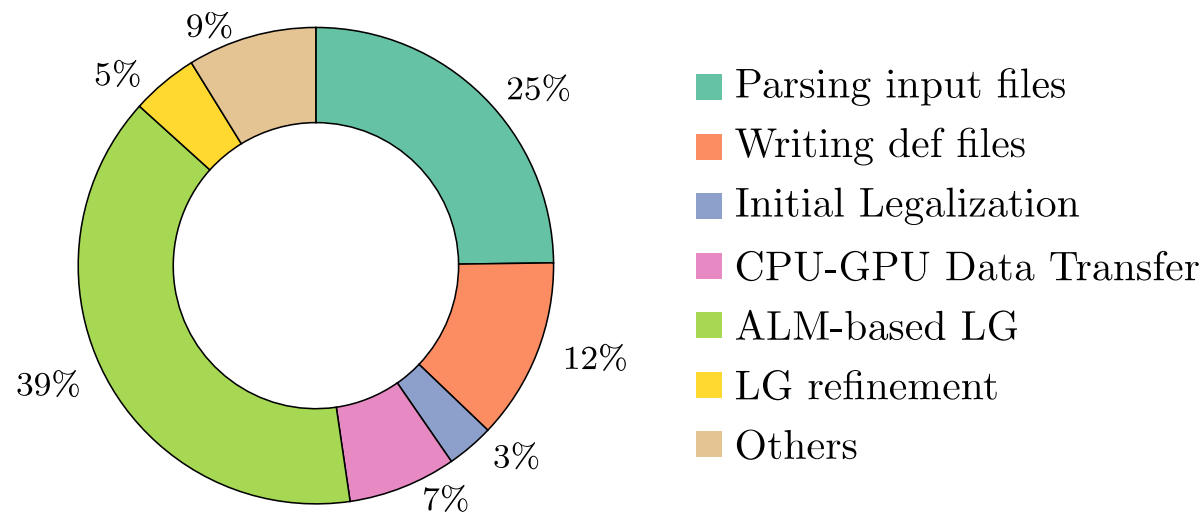
□ Grid Size: Larger sub-regions reduce parallelism but improve movement.



LEGALM-2H: double sub-region size
LEGALM-4H: quadruple sub-region size

□ Runtime Breakdown:

- ALM-based Legalization: 39%.
- Initial Legalization: 3%
- Legalization Refinement: 5%.



Conclusion & Future Work

LEGALM, an efficient inter-row legalization method with the linearized augmented Lagrangian method that supports

- mixed-cell-height circuits
- fence regions

□ Conclusion:

- Linearized ALM formulation for mixed-cell-height legalization.
- Block Gradient Descent (BGD) + Triple-fold partitioning for GPU acceleration.
- 6–36% better quality, $2.25\text{--}5.99\times$ speedup.

□ Future Work:

- Advanced parallelization strategies.

THANK YOU!

 jingmai@pku.edu.cn

Personal Website



magic3007.github.io

